

Der Compiler: Eine Einführung für den Amateur

Berliner Linux User Group

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Outline

Einführung

Lexer

Parser

IR

CodeGen

Optimisierungen

Assembler

Linker

Ressourcen

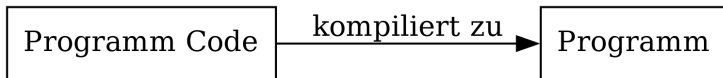
Wer bin ich?

- ▶ Benjamin Stürz <benni@stuerz.xyz>
- ▶ Hobby Compiler Entwickler
- ▶ Aus Berlin
- ▶ Studiere Informatik an der TH Brandenburg
- ▶ OpenBSD User
- ▶ <https://stuerz.xyz/Compiler.pdf>

Die Sprache

```
fn main() {  
    let x = 0;  
    while x < 10 {  
        print x;  
        x = x + 1;  
    }  
    return 0;  
}
```

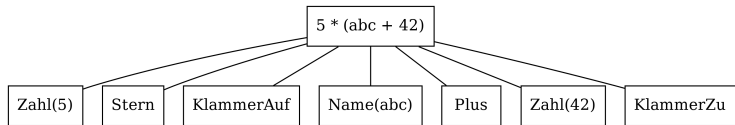
Was ist ein Compiler?



Lexer - Theorie



Lexer - Tokens



Lexer - Regeln I

Zahl	<code>[0-9]+</code>
Name	<code>[a-zA-Z_][a-zA-Z_0-9]*</code>

Lexer - Regeln II

Plus	"+"
Minus	"_"
Stern	"*"
Slash	"/"
Gleich	"="
GleichGleich	"=="
KlammerAuf	" ("
KlammerZu	") "
GeschwKlammerAuf	"{"
GeschwKlammerZu	"}"
Kleiner	"<"
KleinerGleich	"<="
Größer	">"
GrößerGleich	">="
Nicht	"!"
NichtGleich	"!="
Semikolon	";"

Lexer - Regeln III

Return	"return"
Let	"let"
Print	"print"
Fn	"fn"
If	"if"
Else	"else"
While	"while"

Lexer - Implementation - Token

```
enum token_type {
    TK_Zahl,           // Zahl(x)
    TK_Name,           // Name(x)

    TK_Plus,           // "+"
    TK_Minus,          // "-"
    TK_KlammerAuf,     // "("
    TK_KlammerZu,      // ")"
    // Restliche Operatoren...

    TK_Print,          // "print"
    TK_Return,         // "return"
    // Restliche Schlüsselwörter...

    TK_EOF,            // End of File
};
```

```
struct token {
    enum token_type type;
    union {
        int zahl;
        char *name;
    };
};
```

Lexer - Implementation - Übersicht

```
struct token lex(void) {  
    struct token tk;  
    int ch;  
  
    // Leerzeichen überspringen  
    // Zeichen einlesen  
  
    if (isdigit(ch)) {  
        // Zahl lexen  
    } else if (isalpha(ch) || ch == '_') {  
        // Name oder Schlüsselwort lexen  
    } else switch (ch) {  
        // Operatoren  
    default:  
        // Fehler  
    }  
  
    return tk;  
}
```

Lexer - Implementation - Leerzeichen

```
do {  
    ch = getchar();  
} while (isspace(ch));
```

Lexer - Implementation - Zahlen

```
if (isdigit(ch)) {  
    int wert = 0;  
  
    do {  
        wert = wert * 10 + (ch - '0');  
        ch = getchar();  
    } while (isdigit(ch));  
  
    ungetc(ch, stdin);  
  
    tk.type = TK_Zahl;  
    tk.zahl = wert;  
}
```

Lexer - Implementation - Namen & Schlüsselwörter

```
if (isalpha(ch) || ch == '_') {
    size_t len = 0, cap = 10;
    char *s = malloc(cap + 1);

    do {
        if (len == cap) {
            cap *= 2;
            s = realloc(s, cap + 1);
        }
        s[len++] = ch;
        ch = getchar();
    } while (isalnum(ch) || ch == '_');

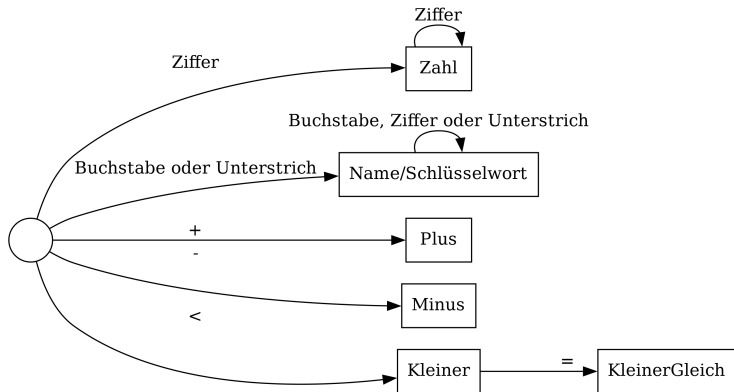
    ungetc(ch, stdin);
    s[len] = '\0';

    if (strcmp(s, "print") == 0) {
        tk.type = TK_Print;
    } else if (strcmp(s, "return") == 0) {
        tk.type = TK_Return;
    } // Weitere Schlüsselwörter...
    else {
        tk.type = TK_Name;
        tk.name = s;
    }
}
```

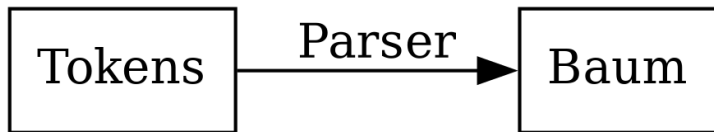
Lexer - Implementation - Operatoren

```
switch (ch) {
case EOF: tk.type = TK_EOF; break;
case '+': tk.type = TK_Plus; break;
case '-': tk.type = TK_Minus; break;
case '*': tk.type = TK_Stern; break;
case '/': tk.type = TK_Slash; break;
case '(': tk.type = TK_KlammerAuf; break;
case ')': tk.type = TK_KlammerZu; break;
case '{': tk.type = TK_GeschwKlammerAuf; break;
case '}': tk.type = TK_GeschwKlammerZu; break;
case ';': tk.type = TK_Semikolon; break;
case '<':
    ch = getchar();
    if (ch == '=') {
        tk.type = TK_KleinerGleich;
    } else {
        tk.type = TK_Kleiner;
        ungetc(ch, stdin);
    }
    break;
// TODO: '>', '>=', '!', '!=', '=', '=='
default:
    errx(1, "invalid input: '%c'", ch);
}
```

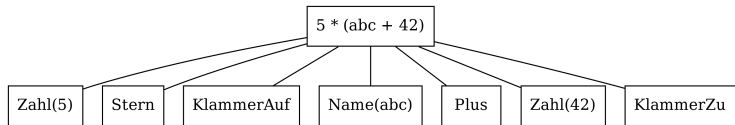
Lexer - Zusammenfassung



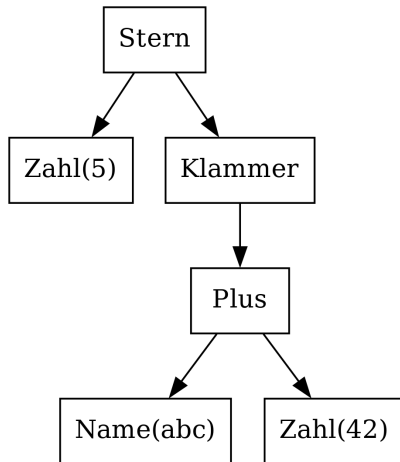
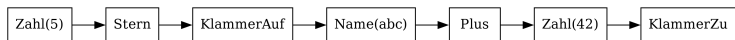
```
// Die einzelnen Typen unserer Token.  
enum token_type;  
// Die Tokens selbst  
struct token;  
// Die Funktion, die uns Tokens zurückgibt.  
struct token lex(void);
```



Parser - Theorie II



Parser - Theorie III



Parser - Regeln I

```
File      = Function+ EOF

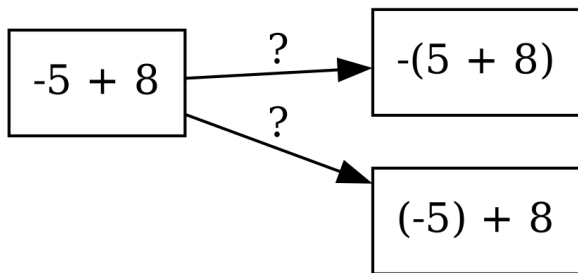
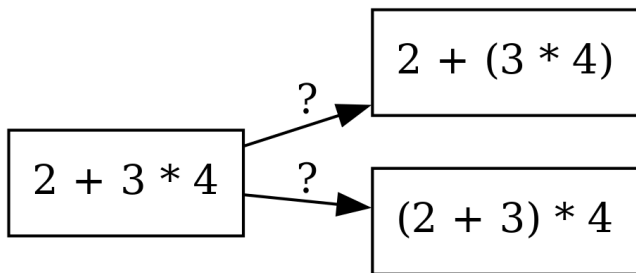
Function  = "fn" Name "(" ")" Block

Block     = "{" Statement* "}"

Statement = "print" Expression ";"
          | "return" Expression ";"
          | "let" Name "=" Expression ";"
          | Name "=" Expression ";"
          | "if" Expression Block ("else" Block)?
          | "while" Expression Block
          | Block

Expression = Zahl
          | Name
          | "(" Expression ")"
          | "-" Expression
          | Expression "+" Expression
          | Expression "-" Expression
          | Expression "*" Expression
          | Expression "/" Expression
          | Expression "==" Expression
          | Expression "!=" Expression
          | Expression "<" Expression
          | Expression "<=" Expression
          | Expression ">" Expression
          | Expression ">=" Expression
```

Parser - Regeln II



Parser - Regeln III

Expression = Comparison

Comparison = Addition "=" Addition
| Addition "!=" Addition
| Addition "<" Addition
| Addition "<=" Addition
| Addition ">" Addition
| Addition ">=" Addition
| Addition

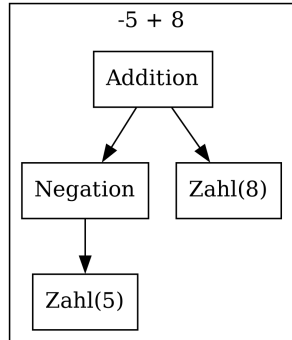
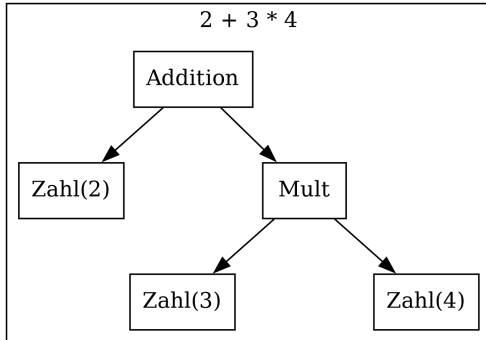
Addition = Addition "+" Mult
| Addition "-" Mult
| Mult

Mult = Mult "*" Unary
| Mult "/" Unary
| Unary

Unary = "-" Unary
| Atom

Atom = Zahl
| Name
| "(" Expression ")"

Parser - Regeln IV



Parser - Regeln V

```
File      = Function+ EOF

Function  = Fn Name KlammerAuf KlammerZu Block

Block     = GeschwKlammerAuf Statement* GeschwKlammerZu

Statement = Print Expression Semikolon
          | Return Expression Semikolon
          | Let Name Gleich Expression Semikolon
          | Name Gleich Expression Semikolon
          | If Expression Block (Else Block)?
          | While Expression Block

Expression = Addition RelOp Addition
          | Addition

RelOp     = GleichGleich | NichtGleich | Kleiner | KleinerGleich | Größer | GrößerGleich

Addition  = Addition (Plus | Minus) Mult
          | Mult

Mult      = Mult (Stern | Slash) Unary
          | Unary

Unary     = Minus Unary
          | Atom

Atom      = Zahl
          | Name
          | KlammerAuf Expression KlammerZu
```

Parser - Datentypen - File, Function, Block

```
struct file {  
    struct function **funcs;  
};  
  
struct function {  
    char *name;  
    struct block *block;  
};  
  
struct block {  
    struct statement **stmts;  
};
```

Parser - Datentypen - Statement

```
enum statement_type {
    ST_Print,
    ST_Return,
    ST_Let,
    ST_Assign,
    ST_If,
    ST_While,
};

struct statement {
    enum statement_type type;
    union {
        struct expression *expr; // print, return
        struct {
            char *name;
            struct expression *value;
        } let;
        struct {
            struct expression *condition;
            struct block *true_case;
            struct block *false_case; // optional
        } ifst; // if statement
        struct {
            struct expression *condition;
            struct block *block;
        } whst; // while statement
    };
};
```

Parser - Datentypen - Expression

```
enum expression_type {
    EX_Zahl,
    EX_Name,
    EX_Klammer,
    EX_Negation,
    EX_Addition,
    EX_Subtraktion,
    EX_Multiplikation,
    EX_Division,
    EX_Gleich,
    EX_Ungleich,
    EX_Kleiner,
    EX_KleinerGleich,
    EX_Größer,
    EX_GrößerGleich,
};

struct expression {
    enum expression_type type;
    union {
        int zahl;
        char *name;
        struct expression *expr; // Klammer, Negation
        struct {
            struct expression *links;
            struct expression *rechts;
        } binär;
    };
};
```

Parser - Implementation - Hilfsfunktionen

```
#define new(T) ((T *)malloc(sizeof(T)))

bool has_token;
struct token peekd;

struct token peek(void) {
    if (!has_token) {
        peekd = lex();
        has_token = true;
    }
    return peekd;
}

struct token next(void) {
    if (has_token) {
        has_token = false;
        return peekd;
    } else {
        return lex();
    }
}

struct token expect(enum token_type type) {
    struct tk = next();
    if (tk.type != type)
        errx(1, "syntax error");
    return tk;
}
```

Parser - Implementation - Atom

```
Atom = Zahl
      | Name
      | KlammerAuf Expression KlammerZu

struct expression *atom(void) {
    struct expression *e = new(struct expression);
    struct token tk;

    switch (tk.type) {
    case TK_Zahl:
        e->type = EX_Zahl;
        e->zahl = tk.zahl;
        break;
    case TK_Name:
        e->type = EX_Name;
        e->name = tk.name;
        break;
    case TK_KlammerAuf:
        e->type = EX_Klammer;
        e->expr = expression();
        expect(TK_KlammerZu);
        break;
    default:
        errx(1, "expected atom");
    }
    return e;
}
```

Parser - Implementation - Unary

```
Unary = Minus Unary  
      | Atom
```

```
struct expression *unary(void) {  
    if (peek().type == TK_Minus) {  
        struct expression *e = new(struct expression);  
        next();  
        e->type = EX_Negation;  
        e->expr = unary();  
        return e;  
    } else {  
        return atom();  
    }  
}
```

Parser - Implementation - Mult

Mult = Mult (Stern | Slash) Unary
| Unary

```
struct expression *mult(void) {  
    struct expression *links = unary();  
  
    while (peek().type == TK_Stern || peek().type == TK_Slash) {  
        struct expression *e;  
  
        if (peek().type == TK_Stern) {  
            e->type = EX_Multiplikation;  
        } else {  
            e->type = EX_Division;  
        }  
        next();  
  
        e->binär.links = links;  
        e->binär.rechts = unary();  
        links = e;  
    }  
  
    return links;  
}
```

Parser - Implementation - Addition

Addition = Addition (Plus | Minus) Mult
 | Mult

```
struct expression *addition(void) {
    struct expression *links = mult();

    while (peek().type == TK_Plus || peek().type == TK_Minus) {
        struct expression *e;

        if (peek().type == TK_Plus) {
            e->type = EX_Addition;
        } else {
            e->type = EX_Subtraktion;
        }
        next();

        e->binär.links = links;
        e->binär.rechts = mult();
        links = e;
    }

    return links;
}
```

Parser - Implementation - Expression

```
RelOp      = GleichGleich | NichtGleich | Kleiner | KleinerGleich | Größer | GrößerGleich
Expression = Addition RelOp Addition
           | Addition

struct expression *expression(void) {
    struct expression *e, *links;
    enum expr_typ type;

    links = addition();

    switch (peek().type) {
        case TK_GleichGleich: type = EX_Gleich;          break;
        case TK_NichtGleich:  type = EX_Ungleich;        break;
        case TK_KleinerGleich: type = EX_KleinerGleich;   break;
        case TK_GrößerGleich: type = EX_GrößerGleich;    break;
        case TK_Kleiner:      type = EX_Kleiner;          break;
        case TK_Größer:       type = EX_Größer;           break;
        default:               return links;
    }
    next();

    e = new(struct expression);
    e->type = type;
    e->binär.links = links;
    e->binär.rechts = addition();
    return e;
}
```

Parser - Implementation - Statement

```
Statement = Print Expression Semikolon  
          | Return Expression Semikolon  
          | Let Name Gleich Expression Semikolon  
          | Name Gleich Expression Semikolon  
          | If Expression Block (Else Block)?  
          | While Expression Block
```

```
struct statement *statement(void) {  
    struct statement *st = new(struct statement);  
    struct token tk = next();  
  
    switch (tk.type) {  
    case TK_Print:  
        st->type = ST_Print;  
        st->expr = expression();  
        expect(TK_Semikolon);  
        break;  
    case TK_If:  
        st->type = ST_If;  
        st->ifst.condition = expression();  
        st->ifst.true_case = block();  
        if (peek().type == TK_Else) {  
            next();  
            st->ifst.false_case = block();  
        } else {  
            st->ifst.false_case = NULL;  
        }  
        break;  
    // ...  
    }  
    return st;  
}
```

Parser - Implementation - Block

```
Block = GeschwKlammerAuf Statement* GeschwKlammerZu
```

```
struct block *block(void) {
    struct block *b = new(struct block);
    size_t len = 0, cap = 10;

    b->stmts = calloc(cap + 1, sizeof(struct statement *));
    expect(TK_GeschwKlammerAuf);

    while (peek().type != TK_GeschwKlammerZu) {
        if (len == cap) {
            cap *= 2;
            b->stmts = reallocarray(b->stmts, cap + 1, sizeof(struct statement *));
        }

        b->stmts[len++] = statement();
    }

    b->stmts[len] = NULL;
    expect(TK_GeschwKlammerZu);
    return b;
}
```

Parser - Implementation - Funktion

Function = Fn Name KlammerAuf KlammerZu Block

```
struct function *function(void) {  
    struct function *f = new(struct function);  
  
    expect(TK_Fn);  
    f->name = expect(TK_Name).name;  
    expect(TK_KlammerAuf);  
    expect(TK_KlammerZu);  
    f->block = block();  
  
    return f;  
}
```

Parser - Implementation - File

```
File = Function+ EOF
```

```
struct file *parse(void) {
    struct file *f = new(struct file);
    size_t len = 0, cap = 10;

    f->funcs = calloc(cap + 1, sizeof(struct function *));

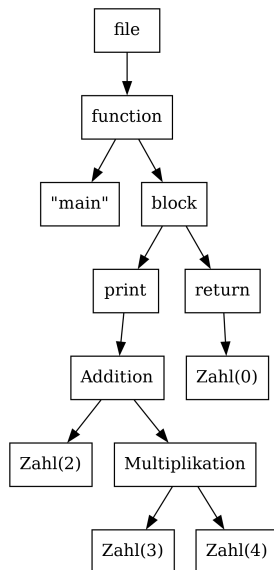
    do {
        if (len == cap) {
            cap *= 2;
            f->funcs = reallocarray(f->funcs, cap + 1, sizeof(struct function *));
        }
        f->funcs[len++] = function();
    } while (peek().type != TK_EOF);

    f->funcs[len] = NULL;

    return f;
}
```

Parser - Zusammenfassung

```
fn main() {  
    print 2 + 3 * 4;  
    return 0;  
}
```



Parser - Was tun mit dem Baum?

- ▶ Ausführen => Interpreter
- ▶ Code Generieren => Compiler
 - ▶ Eine andere Sprache? => Transpiler
 - ▶ Assembler?
 - ▶ IR?

IR - Theorie

- ▶ Intermediate Representation
- ▶ liegt irgendwo zwischen Assembler und C
- ▶ Beispiele: LLVM IR, Java Bytecode, ...

IR - Warum?

- ▶ Einfachere Portierbarkeit auf andere Plattformen
- ▶ Extra Optimisierungen

IR - Beispiel

```
fn main() {  
    print 2 + 3 * 4;  
    return 0;  
}
```

```
fn main() {  
    iload 2  
    iload 3  
    iload 4  
    imul  
    iadd  
    iprint  
    iload 0  
    ireturn
```



► Java

```
fn main() {  
    $1 = 2;  
    $2 = 3  
    $3 = 4;  
    $4 = imul $2, $3;  
    $5 = iadd $1, $4;  
    print $5;  
    $6 = 0;  
    ret $6;  
}
```

► GCC
► LLVM

IR - Generation Expressions

```
void gen_expr(const struct expr *e) {  
    switch (e->type) {  
        case EX_Zahl:  
            printf("iload %s\n", e->zahl);  
            break;  
        case EX_Klammer:  
            gen_expr(e->expr);  
            break;  
        case EX_Addition:  
            gen_expr(e->binär.links);  
            gen_expr(e->binär.rechte);  
            printf("iadd\n");  
            break;  
        // ...  
    }  
}
```

IR - Generation Labels

```
if condition {  
    print 42;  
}
```

wird zu

```
// Code für condition  
  
jmp-if-not .L1  
iload 42  
iprint  
  
.L1:  
// Weiterer Code
```

CodeGen

```
fn main() {  
    print 42;  
}
```

=>

```
section .rodata  
fmt: db "Zahl: %d", 10, 0  
  
section .text  
default rel  
global main  
extern printf  
main:  
    push rbp  
    lea rdi, [fmt]  
    mov rsi, 42  
    xor eax, eax  
    call printf  
    pop rbp  
    ret
```

Optimisierungen - Einführung

TODO: Schnelle Grafik

Optimierungen - Peephole Optimizations

```
$0  ...;  
$1  2;  
$2  imul $0, $1;
```

=>

```
$0  ...;  
$1  1;  
$2  ishl $0, $1;
```

Optimierungen - Dead Code Elimination

```
if 1 == 1 {  
    print 1;  
} else {  
    print 0;  
}
```

=>

```
if 1 == 1 {  
    print 1;  
}
```

Optimierungen - Constant Folding

```
print x + 3 * 4;
```

=>

```
print x + 12;
```

Optimierungen - Viel mehr...

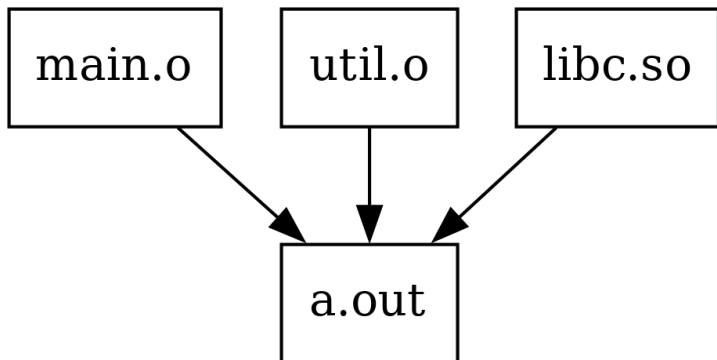
V • T • E		Compiler optimizations	[hide]
Basic block		Peephole optimization • Local value numbering	
Loop		Automatic parallelization • Automatic vectorization • Induction variable • Loop fusion • Loop-invariant code motion • Loop inversion • Loop interchange • Loop nest optimization • Loop splitting • Loop unrolling • Loop unswitching • Software pipelining • Strength reduction	
Data-flow analysis		Available expression • Common subexpression elimination • Constant folding • Dead store elimination • Induction variable recognition and elimination • Live-variable analysis • Use-define chain	
SSA-based		Global value numbering • Sparse conditional constant propagation	
Code generation		Instruction scheduling • Instruction selection • Register allocation • Rematerialization	
Functional		Deforestation • Tail-call elimination	
Global		Interprocedural optimization	
Other		Bounds-checking elimination • Compile-time function execution • Dead-code elimination • Expression templates • Inline expansion • Jump threading • Partial evaluation • Profile-guided optimization	
Static analysis		Alias analysis • Array-access analysis • Control-flow analysis • Data-flow analysis • Dependence analysis • Escape analysis • Pointer analysis • Shape analysis • Value range analysis	

Assembler

```
$ nasm -felf64 -o main.o main.asm
```

Linker

```
$ cc -o program main.o util.o
```



Ressourcen

- ▶ Compilerbau Bücher
- ▶ Wikipedia
- ▶ godbolt.org